

# FPE : Fast Polynomial Evaluation

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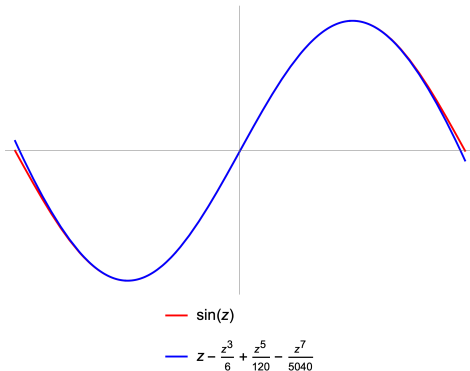
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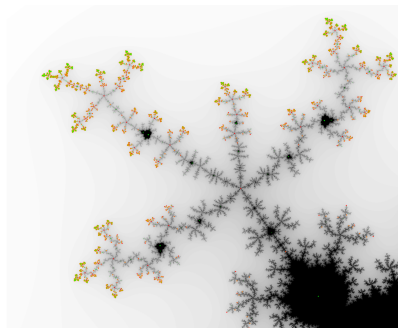
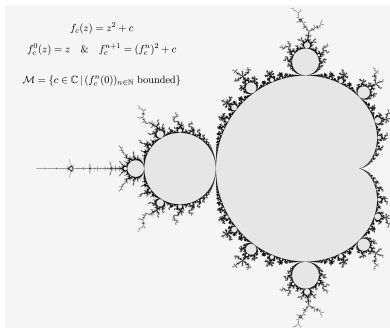
JOURNÉES CALCUL & DONNÉES 2023

# 1. Polynomials

Polynomials are **ubiquitous**:

$$P(z) = a_0 + a_1z + a_2z^2 + \dots + a_dz^d$$





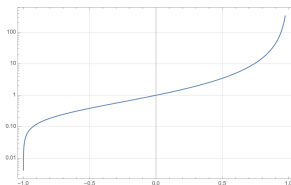
Tera-polynomial (deg =  $10^{12}$ ) slain on Roméo, 2021.

Pre-periodic point (deg =  $10^{10}$ ,  $\times 100$ ) on Roméo, 2023.

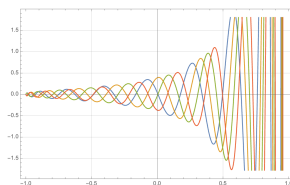
# The Jacobi polynomials $P_n^{(\alpha,\beta)}$

$$\int_{-1}^{+1} P_m^{(\alpha,\beta)}(x) P_n^{(\alpha,\beta)}(x) \cdot (1-x)^\alpha (1+x)^\beta dx = \delta_{m,n}.$$

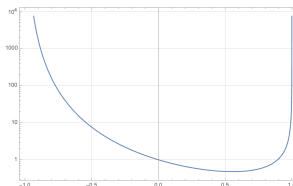
Modulation of a finite energy signal for weight  $\omega_{3,-1}$



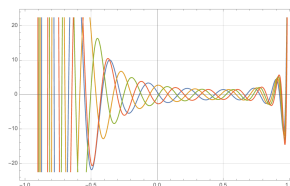
Jacobi polynomials  $P_n^{(3,-1)}(x)$  for  $n=17, \dots, 20$



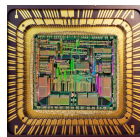
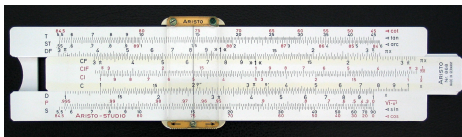
Modulation of a finite energy signal for weight  $\omega_{2,7}$



Jacobi polynomials  $P_n^{(2,7)}(x)$  for  $n=17, \dots, 20$



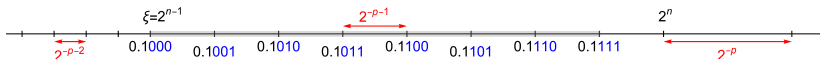
## 2. Floating point arithmetic



Floating point numbers (in base 2) with precision  $p$  bits:

$$\xi = \pm 2^n \times 0.1 \underbrace{\xi_1 \dots \xi_p}_{p \text{ bits}} \quad \text{with} \quad \xi_1, \dots, \xi_p \in \{0, 1\}.$$

The **scale** (order of magnitude) of  $\xi$  is  $n = s(\xi) = 1 + \lfloor \log_2 |\xi| \rfloor$ .



Range of numbers with scale  $n$ .

**Cost of multiplication** in precision  $p$ :

$$M(p) = \begin{cases} \leq 10 \text{ cycles} & \text{FP64 (double, } p = 52) \rightsquigarrow \text{hardware} \\ O(p^2) & \text{schoolbook expansion} \\ O(p^{1.585}) & \text{Karatsuba} \\ O(p^{1.465}) & \text{Toom-Cook} \\ O(p \log p \log \log p) & \text{Schönhage-Strassen} \end{cases}$$

$$(Pz^k + R)(Qz^k + M) = RM + ((P + R)(Q + M) - RM - PQ)z^k + PQz^{2k}$$

$$\boxed{R} \boxed{P} \times \boxed{M} \boxed{Q}$$

# 3. Polynomial evaluation

## Classical methods

How to compute efficiently

$$P(z) = a_0 + a_1z + a_2z^2 + \dots + a_dz^d.$$

- 1 **Monomial:** compute  $z^d$  recursively in time  $O(M(p) \log d)$

$$z^{10} = z^{2+2^3} = (z \times (z^2)^2)^2 \quad z \mapsto z^2 \mapsto z^4 \mapsto z^5 \mapsto z^{10}$$

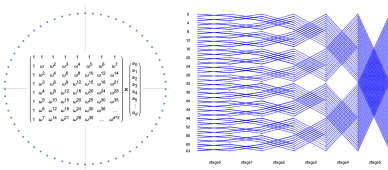
- 2 **Hörner's method:** compute  $P(z)$  in time  $O(M(p)d)$

$$2 - 3z + 4z^2 + 5z^3 = 2 + z(-3 + z(4 + 5z))$$

$$P(z) = a_0 + z \underbrace{(a_1 + z(a_2 + \dots))}_{\text{deg. } d-1}$$

**Multipoint:** compute  $P(z_j)$  on a mesh  $\mathcal{Z} = (z_j)_{0 \leq j \leq d}$ .

- ③ **FFT method:**  $O(M(p)d \log d)$  if  $\mathcal{Z}$  is a regular circular mesh



- ④ **Algebraic methods:**  $O(M(p)d \log^2 d)$  on general mesh, with memory  $O(d^2 \log d)$ .





### 3. Polynomial evaluation

#### Lazy addition

In finite precision ( $p = 4$  decimals), scale considerations can spare us the actual computation:

$$\begin{array}{r}
 z_L \simeq_4 1.234 \times 10^{+2} = 12.340000 \\
 z_S \simeq_4 5.678 \times 10^{-3} = 0.005678 \\
 \hline
 + 12.345678 \simeq_4 1.234 \times 10^{+2}
 \end{array}$$

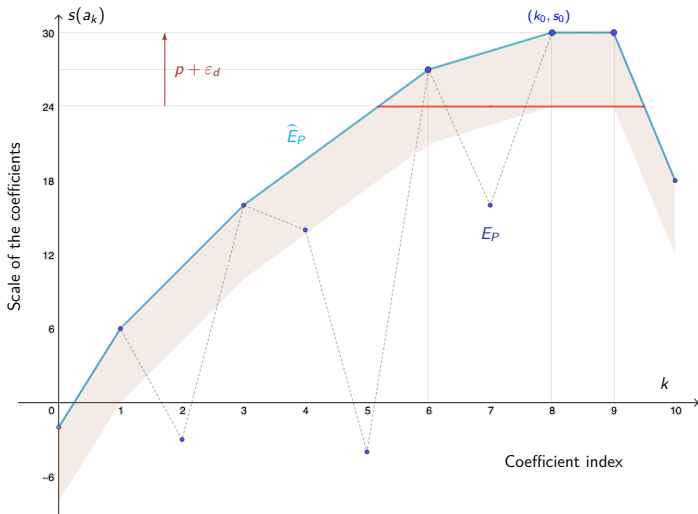
Practical rule: **start with the biggest monomial** and stop when the next operation has no possible influence on the result.

Problems with unsorted coefficients:

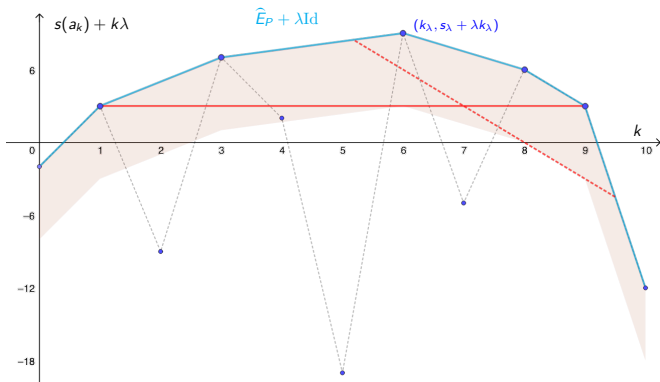
- Lazy evaluation requires sorting  $d$  numbers and costs  $O(d \log d)$  operations  $\gg O(d)$ .
- Absolute precision threshold ?  $s(z_L) \pm 10^{-p} = 0.001 \times 10^{+2}$  requires knowledge of the biggest input.

# FPE algorithm (Fast Polynomial Evaluation)

$$\text{if } |z| = 1, \quad \sum a_k z^k \simeq a_6 + a_8 + a_9$$

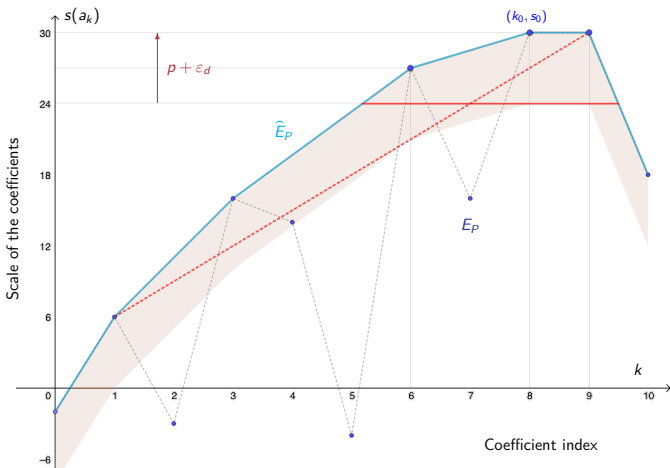


For  $\log_2 |z| = \lambda$ ,  $\sum a_k z^k \simeq a_1 z + a_3 z^3 + a_6 z^6 + a_8 z^8 + a_9 z^9$



It is possible to read the geometric informations on the original graph.

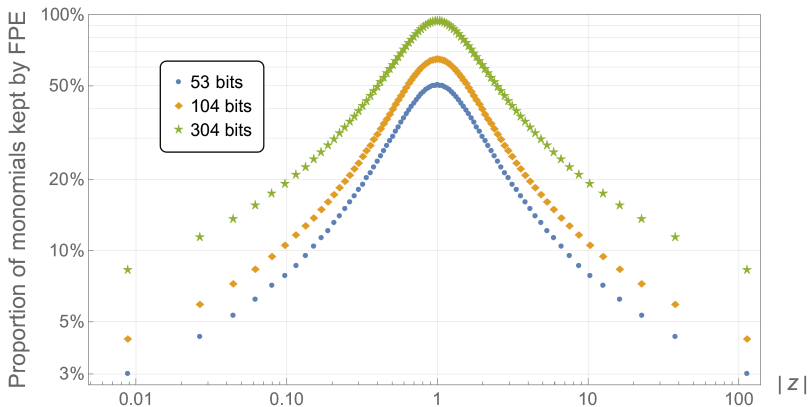
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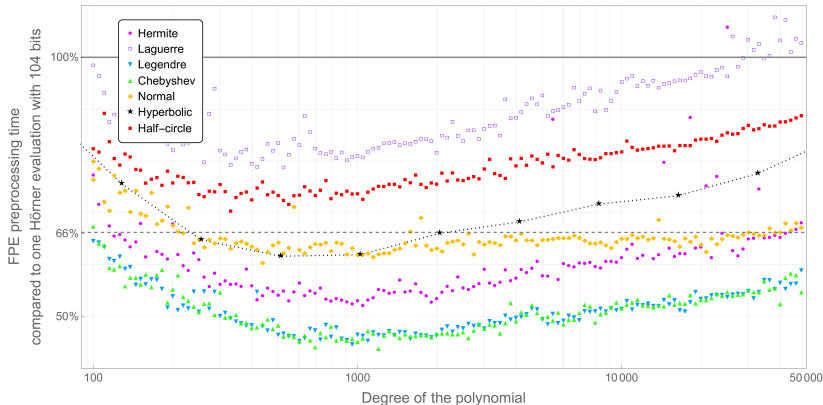
Reduction theorem (based on average width of a thin concave cap):

$$\sum_{k=0}^d a_k z^k \simeq_p \underbrace{\sum_{k \in K_p(z)} a_k z^k}_{\text{Hörner}} \quad \text{with} \quad |K_p(z)| = \underbrace{O\left(\sqrt{d(p + \log d)}\right)}_{\text{on } z\text{-average}}$$

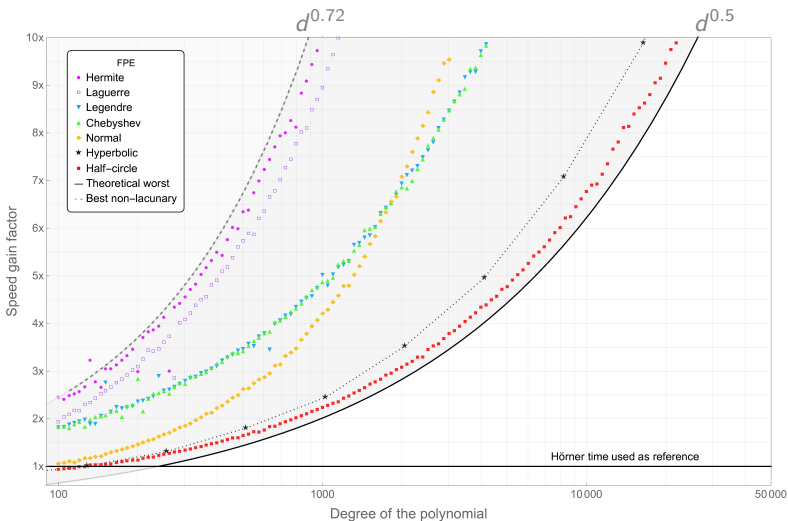


# Benchmark on Roméo

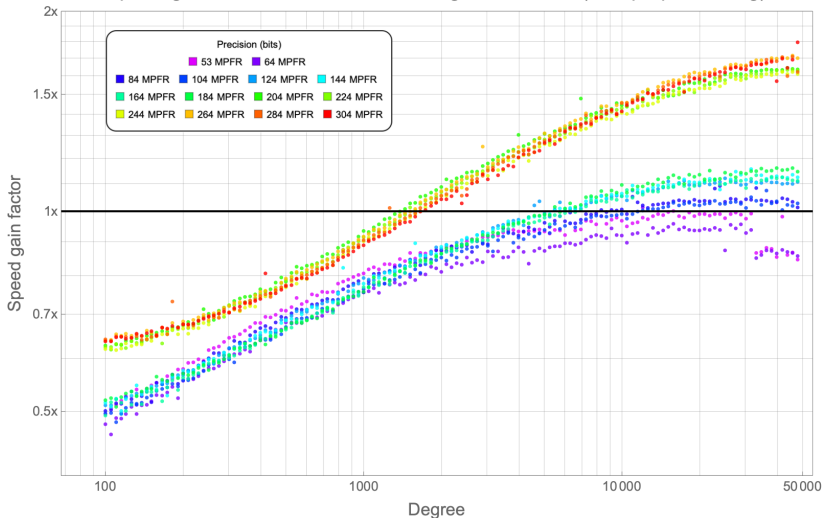
Preprocessing in  $d + 2d \log d$  using only scales (independent of  $p$ ).



Average cost  $\leq M(p)\sqrt{d(p + \log d)}$  per eval. & embarrassingly parallel.

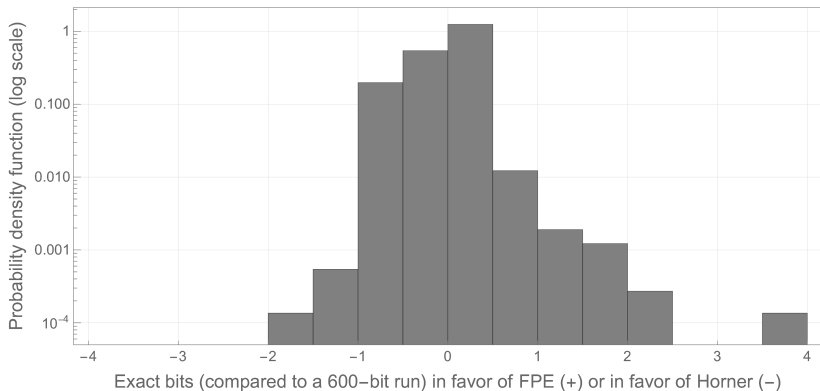


## Speed gain of FPE over Hörner in single evaluation (with preprocessing)

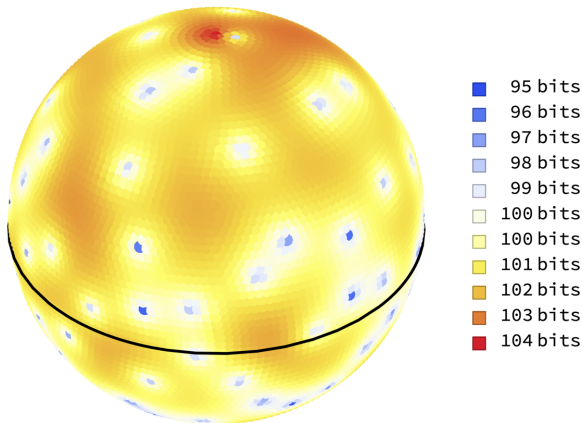




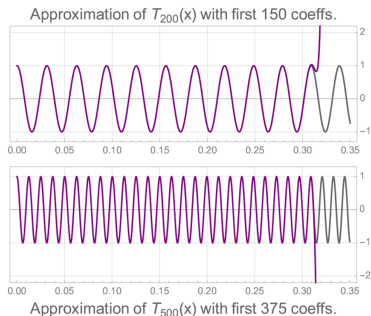
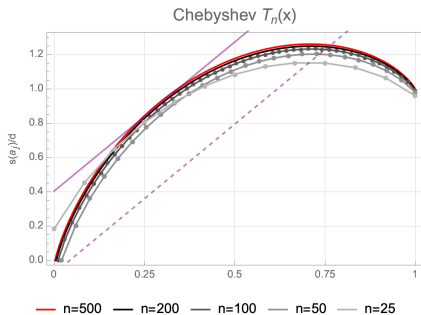
Result guaranteed to be adjacent to the exact value, up to cancelled bits.



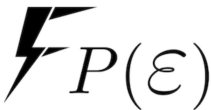
The number of canceled bits (precision loss) is free and explicit.



Predict which terms are necessary.



↪ universal law for Jacobi polynomials (work with Clémentine Courtes).



# FastPolyEval 1.0

Fast Evaluation of Real and Complex Polynomials



<https://github.com/fvigneron/FastPolyEval>

```
git clone fvigneron/FastPolyEval.FPE
make fpe
[vigneron@Francois-MacBook-Pro ~]# ./bin/FastPolyEval -help
$ ./bin/FastPolyEval -help

The following tasks can be performed with the corresponding command line arguments below:
```

Task	Parameters	Short description
-hyperbolic	prec period outFile	writes the coefficients of the hyperbolic polynomial to a CSV file
-Chebyshev	prec degree outFile	writes the coefficients of the Chebyshev polynomial to a CSV file
-Legendre	prec degree outFile	writes the coefficients of the Legendre polynomial to a CSV file
-Hermite	prec degree outFile	writes the coefficients of the Hermite polynomial to a CSV file
-Laguerre	prec degree outFile	writes the coefficients of the Laguerre polynomial to a CSV file
-sum	prec inFile1 inFile2 outFile	computes the sum of two polynomials and writes the result to a CSV file
-diff	prec inFile1 inFile2 outFile	computes the difference of two polynomials and writes the result to a CSV file
-prod	prec inFile1 inFile2 outFile	computes the product of two polynomials and writes the result to a CSV file
-cat	prec inFile1 inFile2 outFile	concatenates two CSV files containing complex numbers
-join	prec inFile1 inFile2 outFile	joins the real part of two sequences into one sequence of complex numbers
-grid	prec inFile1 inFile2 outFile	computes the set product of the real parts of two sequences
-exp	prec inFile outFile	computes the complex exponential of a list of points and writes the result to a CSV file
-rot	prec inFile outFile	maps complex numbers (a, b) to $a+bxp(i)$
-polar	prec inFile outFile	computes the points given by polar coordinates on the sphere and writes the result to a CSV file
-roots	prec inFile outFile	computes the polynomial with a given list of roots and writes the result to a CSV file
-der	prec inFile outFile	computes the derivative of a polynomials and writes the result to a CSV file
-unif	prec count outFile [start] [end]	writes count real numbers in arithmetic progression to a CSV file
-rand	prec count outFile [start] [end]	writes count real random numbers uniformly distributed in an interval to a CSV file
-normal	prec count outFile [center] [var]	writes count real random numbers with Gaussian distribution to a CSV file
-sphere	prec count outFile	writes polar coordinates approximating an uniform distribution on the sphere to a CSV file
-eval	prec polyfile pointfile outFile [checkFile] [count]	quickly evaluates a polynomial on a set of points
-evald	prec polyfile pointfile outFile [checkFile] [count]	quickly evaluates the derivative of a polynomial on a set of points
-evalN	prec polyfile pointfile outFile [checkFile] [count]	quickly evaluates one Newton step of a polynomial on a set of points
-iterN	prec polyfile pointfile outFile [checkFile] [count]	quickly iterates and searches for limits of the Newton method (roots of the polynomial)
-comp	prec inFile1 inFile2 [outfile]	compares two lists of points
-re	prec inFile outFile	writes the real part of the list of complex numbers
-im	prec inFile outFile	writes the imaginary part of the list of complex numbers
-conj	prec inFile outFile	writes the conjugates of the list of complex numbers
-tensor	prec inFile1 inFile2 outFile	computes the tensorial product of the two lists of numbers ( $c_i = a_i * b_i$ )
-analyse	prec polyfile outFile1 [outfile2]	computes the concave cover and the intervals of  z  for which the evaluation strategy changes

For each of the tasks above, use `-task -help` for more detailed information.  
 If the precision is at most 53, machine floating numbers are used (FP64, double),  
 otherwise mpfr floating numbers are used.

# Thank you

